

Contract-based Discovery and Adaptation of Web Services

Luca Padovani

Jointly with
Samuele Carpineti, Giuseppe Castagna, Nils Gesbert,
Cosimo Laneve

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Contracts for Web services

We've got...

- behavioral descriptions of Web services (wsdl, wscl, ws-bpel, put your favorite technology here)
- repositories of Web services descriptions (uddi)

We'd like to...

- look for Web services with a given behavior
- see if it's safe to replace a service with another one
- if not, see whether we can *adapt* one service to safely replace another

We're gonna use...

- **contracts** = abstractions of Web services' behavior

Finding Web services by contract

Compliance = client's satisfaction

$$\rho \dashv \sigma$$

Running a query with *compliance*

$$\mathcal{Q}(\rho) = \{\sigma \mid \rho \dashv \sigma\}$$

Running a query with *duality* ρ^\perp and *subcontract* $\sigma \preceq \tau$

$$\mathcal{Q}(\rho) = \{\sigma \mid \rho^\perp \preceq \sigma\}$$

The quest for \preceq

Desired properties of \preceq

- **reduction** of nondeterminism ($a \oplus b \preceq a$)
- **extension** of functionalities ($a \preceq a + b$)
- some **permutation** of messages ($a.c \preceq c.a$)

The problem

- *reduction* alone is **too strict**
- *extension* is **unsafe**
- *extension;reduction* is **not transitive**
- *permutation* is **not allowed**

The solution

- use (simple) **orchestrators**

Summary

- 1 contracts
- 2 simple orchestrators
- 3 simple orchestrators with buffers
- 4 duality
- 5 recursive behaviors
- 6 orchestrator synthesis
- 7 related and ongoing work

Oh no! More ws-bpe1!

```
<process>
  <sequence>
    <receive operation="Order" variable="Request"/>

    <flow>
      <invoke operation="Deposit" inputVariable="Request" outputVariable="Deposit"/>
      <invoke operation="Charge" inputVariable="Request" outputVariable="Charge"/>
    </flow>

    <switch>
      <case condition="getVariableData(Deposit) == true && getVariableData(Charge) == true">
        <invoke operation="Ship" inputVariable="Request"/>
        <reply operation="Order" value="OK"/>
      </case>

      <case condition="getVariableData(Charge) == true">
        <invoke operation="Refund" inputVariable="Request"/>
        <reply operation="Order" value="NO"/>
      </case>

      <otherwise>
        <reply operation="Order" value="NO"/>
      </otherwise>
    </switch>
  </sequence>
</process>
```

What's in a contract

Actions

O Order
D Deposit
C Charge
S Ship
R Refund

Traces

{ODCSO, OCDSO, ODCRO, OCDRO, ODCO, OCDO}

Branching points

- OD... and OC... is an *external choice*
- ...SO, ...RO, and ...O is an *internal choice*

Basic theory of contracts

Contracts

Syntax

$\sigma ::=$	contract	$\alpha ::=$	action
$\mathbf{0}$	(null)	a	(input)
$\alpha.\sigma$	(action prefix)	\bar{a}	(output)
$\sigma + \sigma$	(external choice)		
$\sigma \oplus \sigma$	(internal choice)		

Example

$0.(\bar{D}.\bar{C}.D.C.(\bar{S}.\bar{0} \oplus \bar{R}.\bar{0} \oplus \bar{0})) + \bar{C}.\bar{D}.D.C.(\bar{S}.\bar{0} \oplus \bar{R}.\bar{0} \oplus \bar{0}))$

Operational semantics

$$\alpha.\sigma \xrightarrow{\alpha} \sigma \quad \sigma \oplus \tau \rightarrow \sigma \quad \frac{\sigma \xrightarrow{\alpha} \sigma'}{\sigma + \tau \xrightarrow{\alpha} \sigma'} \quad \frac{\sigma \rightarrow \sigma'}{\sigma + \tau \rightarrow \sigma' + \tau}$$

Compliance, formally

Systems

$$\rho \parallel \sigma$$

System transitions

$$\frac{\rho \rightarrow \rho'}{\rho \parallel \sigma \rightarrow \rho' \parallel \sigma}$$

$$\frac{\sigma \rightarrow \sigma'}{\rho \parallel \sigma \rightarrow \rho \parallel \sigma'}$$

$$\frac{\rho \xrightarrow{\bar{\alpha}} \rho' \quad \sigma \xrightarrow{\alpha} \sigma'}{\rho \parallel \sigma \rightarrow \rho' \parallel \sigma'}$$

Compliance

$$\rho \dashv \sigma \stackrel{\text{def}}{\iff} \rho \parallel \sigma \implies \rho' \parallel \sigma' \dashv \implies \rho' \xrightarrow{e}$$

Compliance, examples

$$\bar{a}.e \dashv ? a \quad \text{😊}$$

$$\bar{a}.e \oplus \bar{b}.e \dashv ? a + b \quad \text{😊}$$

$$e \dashv ? \sigma \quad \text{😊}$$

$$\bar{a}.e \oplus \bar{b}.e \dashv ? a \oplus b \quad \text{😞}$$

$$\bar{a}.e \dashv ? a \oplus \mathbf{0} \quad \text{😞}$$

$$\mathbf{0} \dashv ? \sigma \quad \text{😞}$$

Subcontract, formally

Set-theoretic interpretation of contracts

$$\llbracket \sigma \rrbracket^s \stackrel{\text{def}}{=} \{ \rho \mid \rho \dashv \sigma \}$$

Subcontract

$$\sigma \sqsubseteq \tau \stackrel{\text{def}}{\iff} \llbracket \sigma \rrbracket^s \subseteq \llbracket \tau \rrbracket^s$$

$$\approx = \sqsubseteq \cap \sqsupseteq$$

Subcontract, (counter)examples

$$\bar{a}.e \quad a + b \not\sqsubseteq a \oplus b$$

$$\bar{a}.e \quad a \not\sqsubseteq \mathbf{0}$$

$$e + a \quad \mathbf{0} \not\sqsubseteq \bar{a}$$

Reduction of nondeterminism

$$\sigma \oplus \tau \sqsubseteq \sigma$$

Properties of strong subcontract

Internal choice = intersection

$$\llbracket \sigma \oplus \tau \rrbracket^s = \llbracket \sigma \rrbracket^s \cap \llbracket \tau \rrbracket^s$$

External choice \neq union

- there are clients in $\llbracket a + b \rrbracket^s$ that are not in $\llbracket a \rrbracket^s \cup \llbracket b \rrbracket^s$:

$$\bar{a}.e \oplus \bar{b}.e \in \llbracket a + b \rrbracket^s \quad \bar{a}.e \oplus \bar{b}.e \notin \llbracket a \rrbracket^s \quad \bar{a}.e \oplus \bar{b}.e \notin \llbracket b \rrbracket^s$$

- sometimes $+$ is \oplus in disguise:

$$\alpha.\sigma + \alpha.\tau \simeq \alpha.(\sigma \oplus \tau)$$

- interferences:

$$\bar{a}.e + \bar{b} \in \llbracket a \rrbracket^s \quad \bar{a}.e + \bar{b} \notin \llbracket a + b \rrbracket^s$$

Properties of strong subcontract

Proposition

\sqsubseteq is a precongurence

$$\sigma \sqsubseteq \tau \quad \Rightarrow \quad \begin{cases} \alpha.\sigma \sqsubseteq \alpha.\tau \\ \sigma \oplus \sigma' \sqsubseteq \tau \oplus \sigma' \\ \sigma + \sigma' \sqsubseteq \tau + \sigma' \end{cases}$$

- + nice axiomatization
- + can be used for safe replacement of **parts** of services

Strong subcontract: axioms

$$(e1) \quad \sigma + \sigma = \sigma$$

$$(e2) \quad \sigma + \tau = \tau + \sigma$$

$$(e3) \quad \sigma + (\sigma' + \sigma'') = (\sigma + \sigma') + \sigma''$$

$$(e4) \quad \sigma + \mathbf{0} = \sigma$$

$$(i1) \quad \sigma \oplus \sigma = \sigma$$

$$(i2) \quad \sigma \oplus \tau = \tau \oplus \sigma$$

$$(i3) \quad \sigma \oplus (\sigma' \oplus \sigma'') = (\sigma \oplus \sigma') \oplus \sigma''$$

$$(d1) \quad \sigma + (\sigma' \oplus \sigma'') = (\sigma + \sigma') \oplus (\sigma + \sigma'')$$

$$(d2) \quad \sigma \oplus (\sigma' + \sigma'') = (\sigma \oplus \sigma') + (\sigma \oplus \sigma'')$$

$$(d3) \quad \alpha.\sigma + \alpha.\tau = \alpha.(\sigma \oplus \tau)$$

$$(d4) \quad \alpha.\sigma \oplus \alpha.\tau = \alpha.(\sigma \oplus \tau)$$

$$(red) \quad \sigma \oplus \tau \leq \sigma$$

Orchestrators

Limitations of \sqsubseteq

\sqsubseteq does not support extensions...

$$a \not\sqsubseteq a + b$$

... because extra actions may cause interferences

$$\bar{a}.e + \bar{b} \dashv\vdash a \qquad \bar{a}.e + \bar{b} \not\sqsubseteq a + b$$

Why not all failures are equal

Failure due to client nondeterminism

$$\bar{a}.e \oplus \bar{b}.e \parallel a \rightarrow \bar{b}.e \parallel a$$

Failure due to service nondeterminism

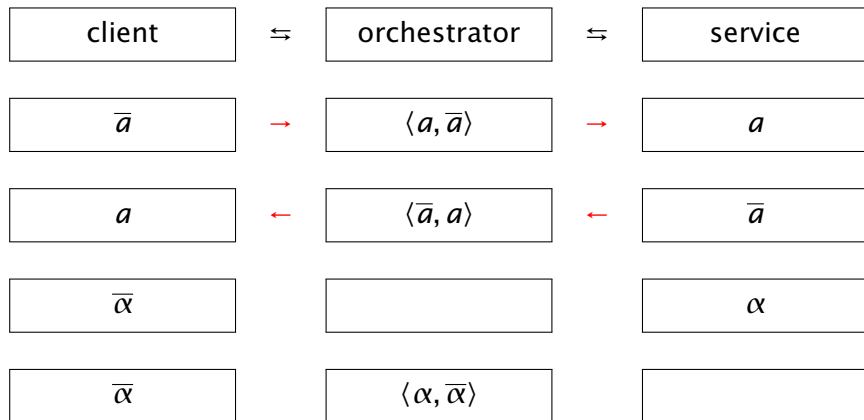
$$\bar{a}.e \parallel a \oplus b \rightarrow \bar{a}.e \parallel a$$

Failure due to “system” nondeterminism

$$\bar{a}.e + \bar{b}.c.e \parallel a + b.\bar{d} \rightarrow c.e \parallel \bar{d}$$

$$\bar{a}.e + \bar{b}.c.e \parallel a + b.\bar{d} \rightarrow e \parallel \mathbf{0}$$

Idea: orchestrated interaction



Synchronous orchestrators

f	::=	orchestrator
		$\mathbf{0}$ (null)
		$\mu.f$ (action prefix)
		$f \vee f$ (disjunction)
μ	::=	orchestration action
		$\langle a, \bar{a} \rangle$ (input/output)
		$\langle \bar{a}, a \rangle$ (output/input)

Orchestrator semantics

Orchestrator transitions

$$\mu.f \xrightarrow{\mu} f \quad \frac{f \xrightarrow{\mu} f'}{f \vee g \xrightarrow{\mu} f'}$$

Trace semantics for orchestrators

$$\llbracket f \rrbracket \stackrel{\text{def}}{=} \{ \mu_1 \cdots \mu_n \mid \exists g : f \xrightarrow{\mu_1} \cdots \xrightarrow{\mu_n} g \}$$

Orchestrated compliance, formally

Orchestrated systems

$$\rho \parallel_f \sigma$$

Orchestrated system transitions

$$\frac{\rho \rightarrow \rho'}{\rho \parallel_f \sigma \rightarrow \rho' \parallel_f \sigma} \quad \frac{\sigma \rightarrow \sigma'}{\rho \parallel_f \sigma \rightarrow \rho \parallel_f \sigma'}$$

$$\frac{\rho \xrightarrow{\bar{\alpha}} \rho' \quad f \xrightarrow{\langle \alpha, \bar{\alpha} \rangle} f' \quad \sigma \xrightarrow{\alpha} \sigma'}{\rho \parallel_f \sigma \rightarrow \rho' \parallel_{f'} \sigma'}$$

Weak compliance

$$f : \rho \dashv\vdash \sigma \stackrel{\text{def}}{\iff} \rho \parallel_f \sigma \implies \rho' \parallel_{f'} \sigma' \dashv\vdash \text{implies } \rho' \xrightarrow{e}$$

Weak compliance, examples

$$\langle a, \bar{a} \rangle : \quad \bar{a}.e \dashv\vdash ? a + b \quad \text{😊}$$

$$\langle a, \bar{a} \rangle \vee \langle b, \bar{b} \rangle : \quad \bar{a}.e \dashv\vdash ? a + b \quad \text{😊}$$

$$\langle a, \bar{a} \rangle \vee \langle b, \bar{b} \rangle : \quad \bar{a}.e \oplus \bar{b}.e \dashv\vdash ? a + b \quad \text{😊}$$

$$\mathbf{0} : \quad e + a \dashv\vdash ? \bar{a} \quad \text{😊}$$

$$f : \quad \bar{a}.e \oplus \bar{b}.e \dashv\vdash ? a \oplus b \quad \text{😞}$$

Weak subcontract, formally

Set-theoretic interpretation of contracts

$$\llbracket \sigma \rrbracket^w \stackrel{\text{def}}{=} \{ \rho \mid \exists f : f : \rho \dashv\vdash \sigma \}$$

Weak subcontract

$$\sigma \preceq \tau \stackrel{\text{def}}{\iff} \llbracket \sigma \rrbracket^s \subseteq \llbracket \tau \rrbracket^w$$

A few doubts...

- is \preceq the subcontract relation we're looking for?
- is \preceq a preorder?

Universal orchestrators

$\sigma \preceq \tau \iff$ for every $\rho, \rho \dashv \sigma$ implies $f : \rho \dashv \tau$ for some f

Universal orchestrator

$f : \sigma \preceq \tau \stackrel{\text{def}}{\iff}$ for every $\rho, \rho \dashv \sigma$ implies $f : \rho \dashv \tau$

f is the *universal orchestrator* for $\sigma \preceq \tau$

Proposition (existence of universal orchestrator)

$\sigma \preceq \tau$ if and only if $f : \sigma \preceq \tau$ for some orchestrator f

Orchestrators as morphisms

Orchestrator application $f(\sigma)$

f	σ	$f(\sigma)$
$\langle a, \bar{a} \rangle \vee \langle b, \bar{b} \rangle$	$a + b$	$a + b$
$\langle a, \bar{a} \rangle \vee \langle b, \bar{b} \rangle$	$a \oplus b$	$a \oplus b$
$\langle a, \bar{a} \rangle$	$a + b$	a
$\langle a, \bar{a} \rangle$	$a \oplus b$	$a \oplus \mathbf{0}$
$\mathbf{0}$	σ	$\mathbf{0}$

Theorem

$f : \rho \dashv\vdash \sigma$ if and only if $\rho \dashv f(\sigma)$

Corollary

$f : \sigma \leq \tau$ if and only if $\sigma \sqsubseteq f(\tau)$

Weak subcontract, examples

Reduction of nondeterminism (\leq embeds \sqsubseteq)

$$\begin{aligned} \langle a, \bar{a} \rangle \vee \langle b, \bar{b} \rangle &: a \oplus b \leq a \\ \Downarrow \\ a \oplus b &\sqsubseteq (\langle a, \bar{a} \rangle \vee \langle b, \bar{b} \rangle)(a) = a \end{aligned}$$

Width extension

$$\begin{aligned} \langle a, \bar{a} \rangle &: a \leq a + b \\ \Downarrow \\ a &\sqsubseteq \langle a, \bar{a} \rangle(a + b) = a \end{aligned}$$

Depth extension

$$\begin{aligned} \mathbf{0} &: \mathbf{0} \leq \sigma \\ \Downarrow \\ \mathbf{0} &\sqsubseteq \mathbf{0}(\sigma) = \mathbf{0} \end{aligned}$$

On transitivity of \preceq

Is \preceq transitive?

$$f : \sigma \preceq \sigma' \quad g : \sigma' \preceq \tau \quad \stackrel{?}{\Rightarrow} \quad h : \sigma \preceq \tau$$

$$\langle a, \bar{a} \rangle \vee \langle b, \bar{b} \rangle . \langle c, \bar{c} \rangle : a \oplus b.c \preceq a$$

$$\langle a, \bar{a} \rangle : a \preceq a + b.d$$

$$??? : a \oplus b.c \preceq a + b.d$$

Proposition (Orchestrator application is monotone)

$\sigma \sqsubseteq \tau$ implies $f(\sigma) \sqsubseteq f(\tau)$

$$\sigma \sqsubseteq f(\sigma') \quad \sigma' \sqsubseteq g(\tau) \quad \Rightarrow \quad \sigma \sqsubseteq f(g(\tau))$$

\preceq is transitive if $f \circ g$ is an orchestrator

Orchestrator composition and transitivity of \preceq

Orchestrator composition

$$f \wedge g$$

- $f \wedge g$ permits the traces permitted by f **and** by g
- $f \wedge g$ forbids the traces forbidden by either f **or** by g

$$\llbracket f \wedge g \rrbracket \stackrel{\text{def}}{=} \llbracket f \rrbracket \cap \llbracket g \rrbracket$$

Proposition

$$f(g(\sigma)) \simeq (f \wedge g)(\sigma)$$

Towards a deduction system for \preceq

Problem: \preceq is not a precongruence w.r.t. $+$

$$\begin{array}{l} \mathbf{0} \preceq a \\ a.b \preceq a.b \\ a.b + \mathbf{0} \not\preceq a.b + a \simeq a.(b \oplus \mathbf{0}) \end{array}$$

Proposition (distributivity of orchestration application)

- 1 $f(\sigma) + f(\tau) \simeq f(\sigma + \tau)$
- 2 $f(\sigma) \oplus f(\tau) \simeq f(\sigma \oplus \tau)$

Corollary

$f : \sigma_1 \preceq \tau_1$ and $f : \sigma_2 \preceq \tau_2$ implies $f : \sigma_1 + \sigma_2 \preceq \tau_1 + \tau_2$

- precongruence is granted if the orchestrator is oblivious of the particular branch taken

Deduction system for \preceq

(red)	$I(\sigma) : \sigma \oplus \tau \leq \sigma$	(prefix)	$\frac{f : \sigma \leq \tau}{\langle \alpha, \bar{\alpha} \rangle . f : \alpha . \sigma \leq \alpha . \tau}$
(width)	$\frac{\llbracket I(\sigma) \wedge I(\tau) \rrbracket = \{\varepsilon\}}{I(\sigma) : \sigma \leq \sigma + \tau}$	(int)	$\frac{f : \sigma \leq \sigma' \quad f : \tau \leq \tau'}{f : \sigma \oplus \tau \leq \sigma' \oplus \tau'}$
(trans)	$\frac{f : \sigma \leq \sigma' \quad g : \sigma' \leq \sigma''}{f \wedge g : \sigma \leq \sigma''}$	(ext)	$\frac{f : \sigma \leq \sigma' \quad f : \tau \leq \tau'}{f : \sigma + \tau \leq \sigma' + \tau'}$

The deduction system is *sound* and *complete*

$$f : \sigma \preceq \tau \iff f : \sigma \leq \tau$$

- completeness regards orchestrators too

Interpretations of orchestrators

As mediators

$$\rho \parallel_f \sigma$$

As morphisms/behavioral coercions

$$f : \sigma \preceq \tau \iff \sigma \sqsubseteq f(\tau) \qquad f : \tau \rightarrow \sigma$$

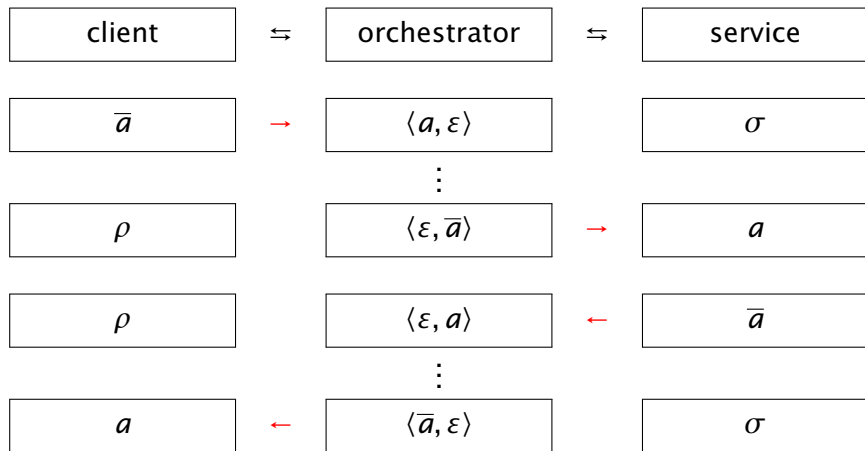
As assumptions on the environment

$$\langle a, \bar{a} \rangle : a \preceq a + b$$

- it is safe to replace a with $a + b$ if nobody ever tries to perform \bar{b}

Buffered orchestrators

Buffered orchestrators



Syntax of buffered orchestrators

$f ::=$	orchestrator
$\mathbf{0}$	(null)
$\mu.f$	(action prefix)
$f \vee f$	(disjunction)
$\mu ::=$	orchestration action
$\langle a, \bar{a} \rangle$	(input/output)
$\langle \bar{a}, a \rangle$	(output/input)
$\langle \alpha, \varepsilon \rangle$	(action/-)
$\langle \varepsilon, \alpha \rangle$	(-/action)

Valid k -orchestrators

Not every orchestrator makes sense

orchestrator	valid	rank
$\langle \varepsilon, a \rangle . \langle \bar{a}, \varepsilon \rangle$	😊	≥ 1
$\langle a, \varepsilon \rangle . \langle a, \varepsilon \rangle$	😊	≥ 2
$\langle \bar{a}, \varepsilon \rangle$	😞	
$\langle \varepsilon, \bar{a} \rangle$	😞	
$\langle a, \varepsilon \rangle . \langle \bar{a}, \varepsilon \rangle$	😞	

Valid k -orchestrators are...

- directional
- finite-state
- fair

Weak k -compliance, formally

Orchestrated systems

$$\rho \parallel_f \sigma$$

Orchestrated system transitions

...

$$\frac{\rho \xrightarrow{\bar{\alpha}} \rho' \quad f \xrightarrow{\langle \alpha, \varepsilon \rangle} f'}{\rho \parallel_f \sigma \rightarrow \rho' \parallel_{f'} \sigma'} \quad \frac{f \xrightarrow{\langle \varepsilon, \bar{\alpha} \rangle} f' \quad \sigma \xrightarrow{\alpha} \sigma'}{\rho \parallel_f \sigma \rightarrow \rho \parallel_{f'} \sigma'}$$

Weak k -compliance

$$f : \rho \dashv\!\! \dashv_k \sigma \stackrel{\text{def}}{\iff} \rho \parallel_f \sigma \implies \rho' \parallel_{f'} \sigma' \dashv\!\! \dashv \text{ implies } \rho' \xrightarrow{e}$$

Weak k -compliance, an example

$$\langle a, \varepsilon \rangle . \langle b, \varepsilon \rangle . \langle \varepsilon, b \rangle . \langle \varepsilon, a \rangle . \langle \bar{c}, c \rangle : \bar{a} . \bar{b} . c . e \dashv\vdash_1 b . a . \bar{c}$$

Weak k -subcontract, formally

Set-theoretic interpretation of contracts

$$\llbracket \sigma \rrbracket_k^w \stackrel{\text{def}}{=} \{ \rho \mid \exists f : f : \rho \dashv\!\! \dashv_k \sigma \}$$

Weak subcontract

$$\sigma \preceq \tau \stackrel{\text{def}}{\iff} \llbracket \sigma \rrbracket^s \subseteq \llbracket \tau \rrbracket_k^w$$

Proposition (existence of universal orchestrator)

$\sigma \preceq_k \tau$ if and only if $f : \sigma \preceq_k \tau$ for some k -orchestrator f

Orchestrators as morphisms

Orchestrator application $f(\sigma)$

f	σ	$f(\sigma)$
$\langle a, \varepsilon \rangle . \langle b, \bar{b} \rangle$	b	$a.b$
$\langle \varepsilon, \bar{a} \rangle . \langle b, \bar{b} \rangle$	$a.b$	b
$\langle a, \varepsilon \rangle . \langle b, \varepsilon \rangle . \langle \varepsilon, \bar{b} \rangle . \langle \varepsilon, \bar{a} \rangle$	$b.a$	$a.b$

Theorem

$f : \rho \dashv|_k \sigma$ if and only if $\rho \dashv f(\sigma)$

Corollary

$f : \sigma \preceq_k \tau$ if and only if $\sigma \sqsubseteq f(\tau)$

Orchestrator composition and transitivity of \preceq_k

We've got a problem

$$\sigma \stackrel{\text{def}}{=} \bar{b} + \bar{d}$$

$$f \stackrel{\text{def}}{=} \langle a, \varepsilon \rangle . \langle c, \varepsilon \rangle . (\langle \varepsilon, \bar{a} \rangle . \langle \bar{b}, b \rangle \vee \langle \varepsilon, \bar{c} \rangle . \langle \bar{d}, d \rangle)$$

$$g \stackrel{\text{def}}{=} \langle a, \varepsilon \rangle . \langle \bar{b}, b \rangle \vee \langle c, \varepsilon \rangle . \langle \bar{d}, d \rangle$$

$$f(g(\sigma)) \simeq f(a.\bar{b} + c.\bar{d}) \simeq a.c.(\bar{b} \oplus \bar{d})$$

No single orchestrator can turn $\bar{b} + \bar{d}$ into $a.c.(\bar{b} \oplus \bar{d})$

Idea

Find an orchestrator $f \cdot g$ such that $f(g(\sigma)) \sqsubseteq (f \cdot g)(\sigma)$

Orchestrator composition, formally

$$\begin{aligned} f \cdot g \stackrel{\text{def}}{=} & \bigvee_{f \xrightarrow{\langle \alpha, \varepsilon \rangle} f'} \langle \alpha, \varepsilon \rangle . (f' \cdot g) \vee \bigvee_{g \xrightarrow{\langle \varepsilon, \bar{\alpha} \rangle} g'} \langle \varepsilon, \bar{\alpha} \rangle . (f \cdot g') \\ & \vee \bigvee_{\langle \varphi, \varphi' \rangle . (f' \cdot g')} \\ & \vee \bigvee_{f \xrightarrow{\langle \varphi, \bar{\alpha} \rangle} f', g \xrightarrow{\langle \alpha, \varphi' \rangle} g', \varphi \varphi' \neq \varepsilon} (f' \cdot g') \\ & \vee \bigvee_{f \xrightarrow{\langle \varepsilon, \bar{\alpha} \rangle} f', g \xrightarrow{\langle \alpha, \varepsilon \rangle} g'} (f' \cdot g') \end{aligned}$$

Theorem

$$f(g(\sigma)) \sqsubseteq (f \cdot g)(\sigma)$$

Deduction system for \preceq_k

No complete deduction system for \preceq_k is known

(swap-inputs)

$$a.b.\sigma = b.a.\sigma$$

(swap-outputs)

$$\bar{a}.\bar{b}.\sigma = \bar{b}.\bar{a}.\sigma$$

(postpone-input)

$$a.\bar{b}.\sigma \leq \bar{b}.a.\sigma$$




Duality

Duality

$$\rho^\perp$$

the \preceq -smallest service that satisfies ρ

Duality, examples

ρ	ρ^\perp
$a.e$	\bar{a}
$a.e \oplus b.e$	$\bar{a} + \bar{b}$
$a.e + b.e$	$\bar{a} \oplus \bar{b}$
e	$\mathbf{0}$
$a.\bar{b}.e + a.\bar{c}.e$	$\bar{a}.(b + c)$
$\mathbf{0}$	
$a.e \oplus \mathbf{0}$	
$a.\mathbf{0}$	
\dots	

Definition (viable contract)

ρ is *viable* if there exists σ such that $\rho \dashv \sigma$

Duality, formally

$$\rho^\perp \stackrel{\text{def}}{=} \sum_{\rho \Downarrow r, e \notin r} \bigoplus_{\alpha \in r, \text{viable}(\rho(\alpha))} \bar{\alpha} \cdot \rho(\alpha)^\perp$$

Theorem

Let ρ be a viable client contract. Then

- 1 $\rho \dashv \rho^\perp$
- 2 $\rho \dashv \sigma$ implies $\rho^\perp \leq_0 \sigma$

Recursive behaviors

Finite syntax for finite behaviors

$\sigma ::=$	contract	$\alpha ::=$	action
	0 (null)		a (input)
	$\alpha.\sigma$ (action prefix)		\bar{a} (output)
	$\sigma + \sigma$ (external choice)		
	$\sigma \oplus \sigma$ (internal choice)		

What about recursive behavior?

Describing recursive behavior

Many different ways...

$$\text{rec } X.a.X \oplus b.\mathbf{0}$$
$$X = a.X \oplus b.\mathbf{0}$$
$$\sigma^*$$

... but it's just syntax!

Infinite syntax for infinite behaviors

$\sigma ::=$	contract	$\alpha ::=$	action
	0		a
	(null)		(input)
	$\alpha.\sigma$		\bar{a}
	(action prefix)		(output)
	$\sigma + \sigma$		
	(external choice)		
	$\sigma \oplus \sigma$		
	(internal choice)		

Definition

The set of contracts is the set of *possibly infinite trees* generated by the grammar above such that

- 1 they have *finitely many different* subtrees
- 2 every *infinite branch* has *infinitely many prefixes*

Every *finite* contracts satisfies these conditions

Examples

$$\begin{aligned} X &= a.X \\ &= a.a.X \\ &= \dots \end{aligned}$$



$$\begin{aligned} X &= a.X \oplus b.0 \\ &= a.(a.X \oplus b.0) \oplus b.0 \\ &= a.(a.(a.X \oplus b.0) \oplus b.0) \oplus b.0 \\ &= a.(a.(a.(a.X \oplus b.0) \oplus b.0) \oplus b.0) \oplus b.0 \\ &= \dots \end{aligned}$$



$$X = X + X$$



$$X = X \oplus X$$



$$X = X$$



Recursion: summary

- all the results stated previously still hold (coinduction)
- use your own preferred syntax (but beware of Kleene *)

Why does it work?

Proposition

Let $D(\sigma) \stackrel{\text{def}}{=} \{\sigma' \mid \sigma \xRightarrow{\varphi} \sigma'\}$. Then $D(\sigma)$ is finite for every σ

Algorithm

Towards an algorithm for deciding \preceq_k

Problem: the orchestrator is not necessarily unique

$$\begin{aligned} \mathbf{0} & : a \oplus b \oplus \mathbf{0} \preceq a + b \\ \langle a, \bar{a} \rangle & : a \oplus b \oplus \mathbf{0} \preceq a + b \\ \langle b, \bar{b} \rangle & : a \oplus b \oplus \mathbf{0} \preceq a + b \\ \langle a, \bar{a} \rangle \vee \langle b, \bar{b} \rangle & : a \oplus b \oplus \mathbf{0} \preceq a + b \end{aligned}$$

$f \preceq g$: g is better (more permissive) than f

$$f \preceq g \stackrel{\text{def}}{\iff} \llbracket f \rrbracket \subseteq \llbracket g \rrbracket$$

Idea

- synthesize the best orchestrator
- the best orchestrator is unique

Deciding \preceq_k

$$\begin{array}{c} \mathbf{a}_r = \{\langle \varphi, \overline{\varphi}' \rangle \mid \sigma \xRightarrow{\varphi}, \tau \xRightarrow{\varphi'}, \mathbb{B} \vdash_k \langle \varphi, \overline{\varphi}' \rangle\} \\ \mathbf{a} = \{\langle \varphi, \overline{\varphi}' \rangle \in \mathbf{a}_r \mid \mathbb{B} \langle \varphi, \overline{\varphi}' \rangle \vdash_k f_{\langle \varphi, \overline{\varphi}' \rangle} : \sigma(\varphi) \preceq \tau(\varphi')\} \\ \tau \Downarrow \mathbf{s} \Rightarrow (\exists r : \sigma \Downarrow r \wedge r \subseteq \mathbf{a} \circ \mathbf{s}) \vee (\emptyset \bullet \mathbf{a}) \cap \overline{\mathbf{s}} \neq \emptyset \\ \hline \mathbb{B} \vdash_k \bigvee_{\mu \in \mathbf{a}} \mu \cdot f_\mu : \sigma \preceq \tau \end{array}$$

Theorem

The following properties hold:

- 1 (termination) the algorithm always terminates
- 2 (correctness) $f : \sigma \preceq_k \tau$ implies that f has rank k and $f : \sigma \preceq_k \tau$
- 3 (completeness) $f : \sigma \preceq_k \tau$ implies $g : \sigma \preceq_k \tau$ for some $g \geq f$

An example from Wil's lecture (1/3)

$$\sigma \stackrel{\text{def}}{=} \text{order} . (\text{money} + \overline{\text{food}} . \text{money})$$

$$\rho_1 \stackrel{\text{def}}{=} \overline{\text{order}} . \text{food} . \overline{\text{money}} . e$$

$$\rho_1^\perp = \text{order} . \overline{\text{food}} . \text{money}$$

$$f_1 = \langle \text{order}, \overline{\text{order}} \rangle . \langle \overline{\text{food}}, \text{food} \rangle . \langle \text{money}, \overline{\text{money}} \rangle$$

An example from Wil's lecture (2/3)

$$\sigma \stackrel{\text{def}}{=} \text{order} . (\text{money} + \overline{\text{food}} . \text{money})$$

$$\rho_2 \stackrel{\text{def}}{=} \overline{\text{order}} . (\text{food} . \overline{\text{money}} . \text{e} + \overline{\text{money}} . \text{food} . \text{e})$$

$$\rho_2^\perp = \text{order} . (\overline{\text{food}} . \text{money} \oplus \text{money} . \overline{\text{food}})$$

$$f_2 = \langle \text{order}, \overline{\text{order}} \rangle . \langle \overline{\text{food}}, \text{food} \rangle . \langle \text{money}, \overline{\text{money}} \rangle$$

An example from Wil's lecture (3/3)

$$\sigma \stackrel{\text{def}}{=} \text{order} . (\text{money} + \overline{\text{food}} . \text{money})$$

$$\rho_3 \stackrel{\text{def}}{=} \overline{\text{order}} . \overline{\text{money}} . \text{food} . e$$

$$\rho_3^\perp = \text{order} . \text{money} . \overline{\text{food}}$$

$$f_3 = \langle \text{order}, \overline{\text{order}} \rangle . \langle \text{money}, \varepsilon \rangle . \langle \overline{\text{food}}, \text{food} \rangle . \langle \varepsilon, \overline{\text{money}} \rangle$$

Conclusions

Wrap-up

Subcontract relation

- tool for *searching* and *reasoning about* services by their **contracts** (= behavioral types)
- \preceq combines **reduction**, **extension**, and **permutation** into a single preorder
- \preceq gives **safe substitution** of services modulo **orchestration**
- \preceq is decidable

(Simple) orchestrators

- have nice properties (**universality**, **compositionality**)
- can be automatically synthesized

Contracts vs session types

What is being typed

- contract = type of a process
- session type = type of a channel
- session type = type of a process projected on a channel

Testing approach

- session types: subtyping *preserves* correctness
- contracts: compliance *defines* subcontracting

Essential bibliography

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