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The Parametric  $\lambda$ -Calculus  
*a Metamodel for Computation*

**ERRATA CORRIGE:** June 7, 2006

Texts in Theoretical Computer Science

An EATCS Series

Springer-Verlag

Berlin Heidelberg New York

London Paris Tokyo

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Budapest

## PART I. Syntax

Section	Original	Correct
1.2 The $\Delta$ -Calculus Page/s: 10 Line/s: 1,2	Thus $KI$ has degree <b>0</b> and $II$ has degree <b>1</b> .	Thus $II$ has degree <b>0</b> and $KI$ has degree <b>1</b> .
1.2.1 Proof of Confluence and Standardization Theorems Page/s: 14 Line/s: 8	$M \hookrightarrow_{\Delta} M', N \hookrightarrow_{\Delta} N'$ and $N \notin \Delta$ imply $MN \hookrightarrow_{\Delta} M'N'$ .	$M \hookrightarrow_{\Delta} M', N \hookrightarrow_{\Delta} N'$ and $MN$ is not a $\Delta$ -redex imply $MN \hookrightarrow_{\Delta} M'N'$ .
2.1.2 Proof of Böhm's Theorem Page/s: 29 Line/s: 6	if $c \equiv \epsilon$ then either $ p - m  \neq  q - n $ or $x \neq y$ ;	if $M =_{\eta} \lambda x_1 \dots x_p. x N_1 \dots N_m$ , $N =_{\eta} \lambda x_1 \dots x_q. y N_1 \dots N_n$ and $c \equiv \epsilon$ then either $ p - m  \neq  q - n $ or $x \neq y$ ;
2.1.2 Proof of Böhm's Theorem Page/s: 29 Line/s: 7	$N =_{\eta} \lambda x_1 \dots x_p. y N_1 \dots N_m$	$N =_{\eta} \lambda x_1 \dots x_p. x N_1 \dots N_m$
2.1.2 Proof of Böhm's Theorem Page/s: 29 Line/s: 15	$U_n^i \equiv \lambda x_1 \dots x_i. x_n$	$U_n^i \equiv \lambda x_1 \dots x_n. x_i$
3.1.2 Proof of Potential $\Gamma$ -Valuability and $\Gamma$ -Solvability Theorem Page/s: 43 Line/s: -2	$M[P/z] \rightarrow_{\Delta} N[Q/z]$	$M[P/z] \rightarrow_{\Delta}^* N[Q/z]$

## PART II. Operational Semantics

Section	Original	Correct
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## PART III. Denotational Semantics

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**PART III. Computational Power**

Section	Original	Correct
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