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The Parametric λ-Calculus

a Metamodel for Computation

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PART I. Syntax

Section	Original	Correct
1.2 The Δ-Calculus Page/s: 10 Line/s: 1,2	Thus KI has degree 0 and II has degree 1.	Thus II has degree 0 and KI has degree 1.
1.2.1 Proof of Confluence and Standardization Theorems Page/s: 14 Line/s: 8	$M \hookrightarrow_{\Delta} M'$,$N \hookrightarrow_{\Delta} N'$ and $N \notin \Delta$ imply $MN \hookrightarrow_{\Delta} M'N'$.	$M \hookrightarrow_{\Delta} M'$,$N \hookrightarrow_{\Delta} N'$ and MN is not a Δ-redex imply $MN \hookrightarrow_{\Delta} M'N'$.
2.1.2 Proof of Böhm's Theorem Page/s: 29 Line/s: 6	if $c \equiv \epsilon$ then either $\vert p-m\vert \neq \vert q-n\vert$ or $x \not\equiv y$;	if $M =_{\eta} \lambda x_1 \ldots x_p.xN_1 \ldots N_m$, $N =_{\eta} \lambda x_1 \ldots x_q.yN_1 \ldots N_n$ and $c \equiv \epsilon$ then either $\vert p-m\vert \neq \vert q-n\vert$ or $x \not\equiv y$;
2.1.2 Proof of Böhm's Theorem Page/s: 29 Line/s: 7	$N =_{\eta} \lambda x_1 \ldots x_p.yN_1 \ldots N_m$	$N =_{\eta} \lambda x_1 \ldots x_p.xN_1 \ldots N_m$
2.1.2 Proof of Böhm's Theorem Page/s: 29 Line/s: 15	$U^i_n \equiv \lambda x_1 ... x_i.x_n$	$U^i_n \equiv \lambda x_1 ... x_n.x_i$
3.1.2 Proof of Potential Γ-Valuability and Γ-Solvability Theorem Page/s: 43 Line/s: -2	$M[P/z] \rightarrow_{\Delta} N[Q/z]$	$M[P/z] \rightarrow^{*}_{\Delta} N[Q/z]$

PART II. Operational Semantics

Section	Original	Correct

PART III. Denotational Semantics

Section	Original	Correct

PART III. Computational Power

Section	Original	Correct